

Efficient Error Estimating Coding: Feasibility and Applications

Binbin Chen, Ziling Zhou, Yuda Zhao & Haifeng Yu

Presented by Qingxi Li

Outline

- Motivation
- Application
- Algorithm
- Related Work

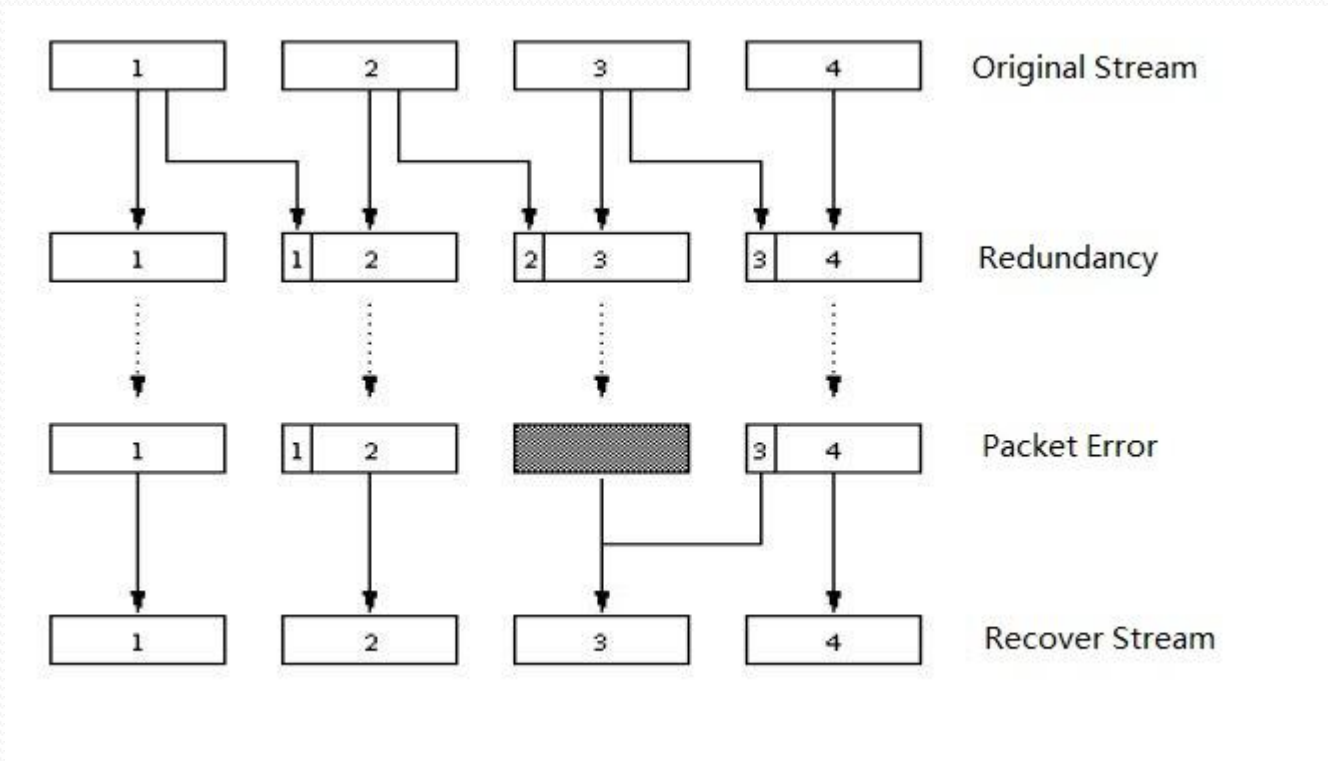
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Motivation

- Flips:
 - 0 ->1 or 1->0
- Traditional solution – drop packet
 - The application can or should use completely correct packets
 - Wastes time and bandwidth
- Partially correct packets can be used
 - Incremental redundancy from the source to recover the packet

Motivation

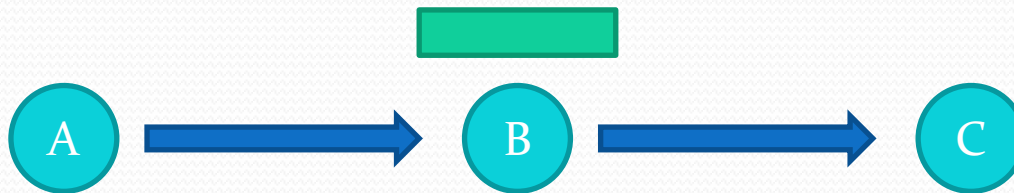


Motivation

- Flips:
 - 0 ->1 or 1->0
- Traditional solution – drop packet
 - The application can or should use completely correct packets
 - Wastes time and bandwidth
- Some partially correct packets can be used
 - Incremental redundancy from the source to recover the packet
 - Directly use the partial packets, like in video

Motivation

- Not all the partially correct packet can be used
 - Determined by Bit Error Rate (BER)
 - Drop if Bit Error Rate (BER) over some threshold
 - Still have bandwidth & time waste
- Goal of this paper: Estimate the bit error rate



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 - Server side
 - Client side
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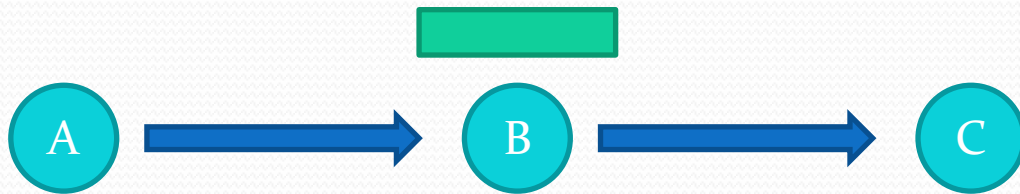
Application

- Senders
 - Transmission rate, channel, power ...
 - Trade off
 - High transmission rate => high throughput & high lost/error rate



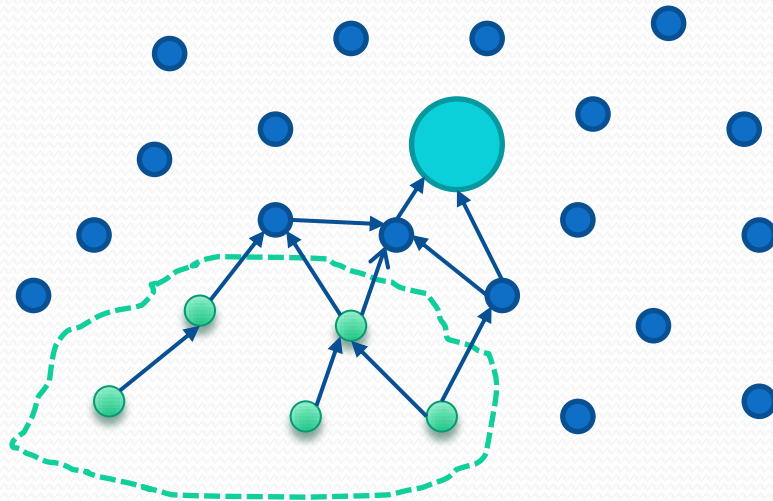
Application

- Client
 - Packet retransmission



Application

- Client
 - Packet retransmission
 - Packet forwarding priority
 - Emergency response in WSN

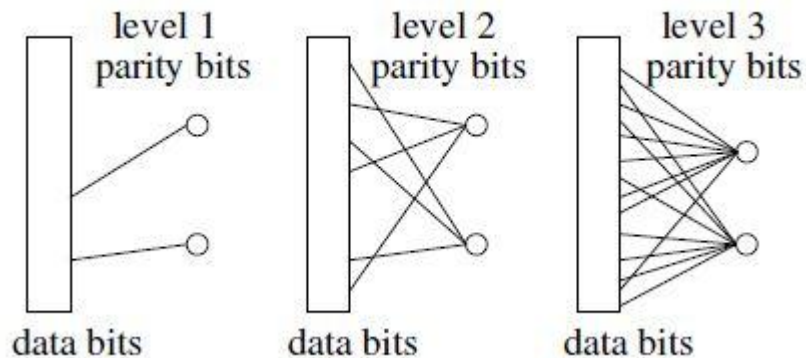


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Algorithm

- Encoding (client side) :
 - $\log_2 n$ levels & s bits.
 - For level i , randomly choose $2^i - 1$ bits to calculate parity bit, repeat s times.
 - Parity bit: even parity bit & odd parity bit



$s=2$

Algorithm

- Estimating (server side) :
 - Assume EBR is in $[a, b]$ ($a=0, b=1/4$)

Algorithm 2 EEC Estimating Procedure.

```
1: for  $i = 1$  to  $i = \lfloor \log_2 n \rfloor$  do
2:   Compute the fraction ( $q_i$ ) of parity bits at level  $i$  that fail
   parity check;
3:   if  $q_1 \geq c_2$  then
4:     Output  $\hat{p} = 1/4$  and exit;
5:   end if
6:   if  $c_1 < q_i < c_2$  then
7:     Output  $\hat{p} = q_i/2^i$  and exit;
8:   end if
9: end for
10: Output  $\hat{p} = 0$  and exit;
```

Algorithm

- Time complexity
 - $s \sum 2^i = O(sn) = O(n)$ (both decoding and encoding)
- Increasing packet size
 - $s \log_2 n = O(\log n)$
 - For some application only want to know whether EBR exceeds a , we only care about the level i s.t. $a \in [2^{-i-1}, 2^{-i}]$
 - 2% of the packet size

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 - Efficient Channel-aware Rate Adaptation in Dynamic Environments
 - Simple Adaptive Relaying Protocol for Wireless Relay Networks

Efficient Channel-aware Rate Adaptation in Dynamic Environments

- Optimal transmission rate depends on environment
 - Higher transmission rate increase throughput & reduce the transmission time
 - ...but reduces the range at which the transmission can be successfully decoded
- How to pick the optimal transmission rate
 - Static channel
 - Trail-and-error
 - Dynamic channel (vehicle, mobile)
 - Properties of the channel change quickly

Efficient Channel-aware Rate Adaptation in Dynamic Environments

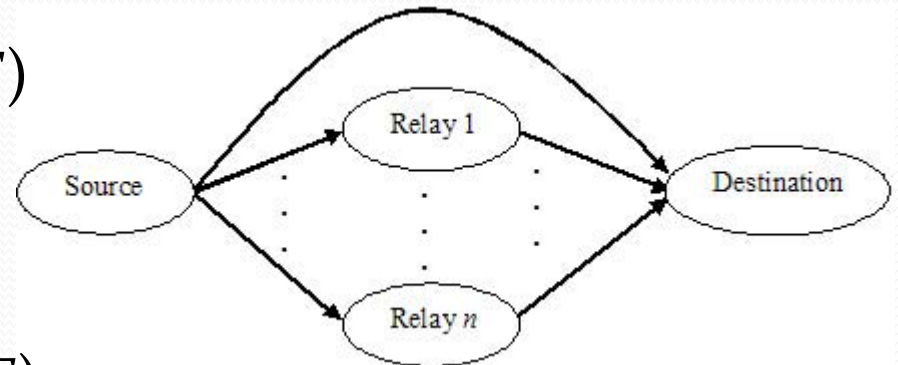
- Successful packet transmission determined by signal-to-interference and noise ratios (SINR)
- SINR determined by
 - received signal strength (RSS)
 - Noise & interference
- Measurement
 - Sender's RSS – RSS indicator
 - Receiver's noise level
 - Receiver's transmission power

Efficient Channel-aware Rate Adaptation in Dynamic Environments

- Estimate path loss
 - Path loss is the energy reduce during transmission
 - path loss = receiver's transmit power – sender's received signal strength
- Predict the path loss in the future
- Estimate SINR at receiver
 - $RSS = \text{transmit power} - \text{path loss}$
 - Noise level measured by receiver
- For each transmission rate, determine a good possible SINR by history of successful/fail transmission
- Choose the transmission rate based on the SINR

Simple Adaptive Relaying Protocol for Wireless Relay Networks

- Wireless relay networks
 - Reduced signal transmit power
- Relaying protocols
 - Decode And Forward(DAF)
 - Reduce the noise
 - error propagation if error rate is large
 - Amplify And Forward(AAF)
 - Amplify the noise & won't suffer error propagation
- Relays can correctly decode the signal will use DAF otherwise AAF





Q & A

Motivation

- Flips:
 - 0 ->1 or 1->0
- Traditional
 - The application can or should use completely correct packets
 - Waster time and bandwidth
- Current situation
 - Partially correct packets can be used
 - Incremental redundancy from the source to recover the packet
 - Directly use the partial packets, like in video